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#### Toward Automatic Performance Tuning for Numerical Simulations in the SILC Matrix Computation Framework

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#### Outline

- The SILC matrix computation framework

   Easy-to-use interface for matrix computation libraries
- Automatic performance tuning in SILC – Performance modeling
  - Related work
- · Experimental results
- · Concluding remarks

### Numerical simulations

- A common feature: Appearance of PDEs
  - Discretization of PDEs results in linear systems, which are solved by linear solvers
- Some simulations require only one linear system to be solved
  - E.g. Steady-state simulations in linear problems
- Others require many linear systems to be solved repeatedly
  - E.g. Steady-state simulations in nonlinear problems, time-dependent simulations

#### Matrix computation libraries

- · Key components of numerical simulations
- · Problem: Using libraries is not easy
  - Many libraries having different APIs
  - Diversity of computing environments
  - Interoperability issues of various programming languages
- Solution: The SILC framework

#### Simple Interface for Library Collections

- Benefits
  - Independent of libraries, environments & languages
  - Easy to use
- · Three steps to use libraries
  - Depositing data (matrices, vectors, etc.) to a server
  - Making requests for computation by means of mathematical expressions
  - Fetching the results of computation if necessary





### Functionalities of SILC

- Data structures for matrix computations – Matrices (dense, band, sparse), vectors, etc.
- Math operators, functions, and subscript
   2-norm of vector x: sqrt(x' \* x)
  - $-5 \times 5$  submatrix of A: A[1:5, k:k+4]
- · No loops and conditional branching
  - These are realized with the languages used to write user programs for SILC

## Main characteristics of SILC

- Independence from programming languages – User programs for SILC in your favorite languages
- Independence from libraries and environments
  - Using alternative libraries and environments requires no modification in user programs
  - Flexible combinations of client & server environments

User program (client)	SILC server
Sequential	Sequential
Sequential	Shared-memory parallel (OpenMP)
Sequential	Distributed parallel (MPI)
Distributed parallel (MPI)	Distributed parallel (MPI)
Distributed parallel (MPI)	Distributed parallel (MPI)

#### Cost in SILC: Communication time

· Likely to be smaller than computation time

Matrix (Solver)	Time for solving a linear system A <b>x</b> = <b>b</b>	Time for depositing A and <b>b</b> and fetching <b>x</b>
Dense (LU decomposition)	O(N <sup>3</sup> )	O(N <sup>2</sup> )
Sparse (CG method)	O(CZ)	<i>O</i> ( <i>Z</i> )

(N: dimension, C: iteration count, Z: number of non-zero elements)

• Possible speedups by parallel computation even at the cost of data communications

#### Automatic performance tuning (APT)

- SILC needs APT
  - To achieve as much speedup as possible in order to relatively minimize the cost of data communications
  - Using all available processors (or threads) is not always optimal
- SILC is an ideal framework in which APT is implemented
  - SILC servers can carry out various types of APT independently of user programs

#### Purposes of the present research

- Performance modeling of time-dependent simulations in SILC
- Outline of an APT mechanism for SILC
- Assumptions
  - A sequential user program, running with
  - A shared-memory parallel SILC server

# Performance modeling

• The execution time (in seconds) of a user program is modeled as a function of *p* (the number of threads) as follows:

f(p) = a/p + bp + c

- a/p: time for parallelized computations
- bp : parallelization overhead
- $\boldsymbol{c}$  : time for sequential computations
- (a, b, c > 0)

#### The least squares method

 Suppose we have measured the execution time of the user program with *n* different numbers of threads (e.g., p<sub>i</sub> = 2<sup>i-1</sup>, i = 1, ..., n)

Number of threads	$p_1$	$p_2$	 $p_n$
Execution time [sec.]	$f_1$	$f_2$	 $f_n$

• By using the least squares method, we can find *a*, *b*, and *c* that minimize

$$q = \sum_{i=1}^{n} \{ f_i - f(p_i) \}$$

#### The optimal number of threads $p_{\rm opt}$

- With a performance model *f*(*p*), we can predict the optimal number of threads *p*<sub>opt</sub> that leads to the minimum execution time
- Since f > 0,  $p_{opt}$  satisfies

$$\frac{df}{dp} = -\frac{a}{p^2} + b = 0$$

• By solving the equation for p, we have  $p = \sqrt{a/b}$  and thus

$$p_{\text{opt}} = \begin{cases} \lfloor p \rfloor & \text{if } f(\lfloor p \rfloor) < f(\lfloor p \rfloor) \\ \lceil p \rceil & \text{otherwise} \end{cases}$$

#### Proposed performance modeling

- 1. Measure the execution time of a user program with different numbers of threads
- 2. Learn a performance model

$$f(p) = a/p + bp + c$$

with the 3 parameters a, b, and c determined by the least squares method

3. Predict the optimal number of threads  $p_{opt}$ 

#### Related work

- Various approaches to detailed performance modeling
- E.g. Performance Analysis and Characterization Environment (PACE) by Kerbyson *et al.* 
  - Semi-automated code analysis of user programs
  - Predefined hardware models
- Our performance modeling is much simpler
   Due to the primary objective of SILC: Independence from libraries, environments, and languages

#### Numerical experiments

- Purpose: Validation of the performance modeling
- · Example applications
  - 1. Cloth simulation based on the implicit Euler method
  - 2. CFD simulation based on the Moving Particle Semi-implicit (MPS) method
  - 3. An initial value problem of the 2-dimensional diffusion equation

#### Test environments

User programs (clients)	Parallel SILC server
Dell Dimension 8400	SGI Altix 3700
Intel Pentium 4 3.4 GHz,	32 Intel Itanium 2 1.3 GHz,
1 GB RAM,	32 GB RAM (cc-NUMA),
Microsoft Windows XP SP2	Red Hat Linux AS 2.1
MinGW (GCC 3.2.3)	Intel C Compiler 9.1
-03 option enabled	-03 option enabled

· Both machines in the same Gigabit Ethernet LAN

#### Validation criterion

 Relative error ε<sub>rel</sub> in the execution time t measured with p<sub>opt</sub>

$$\varepsilon_{\rm rel} = \frac{|t_{\rm true} - t|}{t_{\rm true}}$$

where  $t_{\text{true}}$  is the execution time measured with the true optimal number of threads ( $\varepsilon_{\text{rel}} = 0$  if  $p_{\text{opt}}$  is truly optimal)











/* 1. Calculate $f$ , $\partial f/\partial x$ and $\partial f/\partial y$ */	time steps
<pre>SILC_EXE((") = Y_L * X - Y_K * X ); SILC_EXE((") = sparse(P_row, P_col, p, 3*s, s)"); SILC_EXE(("a = sqrt(diagve((P * P))"); SILC_EXE(("ij = P * (K_stiff *@ (z - L) /@ z)"); SILC_EXE(("a = Y_L * v - Y_R * v"); SILC_EXE(("a = yL * v - Y_R * v"); SILC_EXE(("ij = Q * K_damp"); SILC_EXE(("ij = Q * K_damp"); SILC_EXE(("i = Sum_f * (fij + dij) - M * g");</pre>	
<pre>SILC_EXEC("zhat = ones(s, 1) /@ z; Pzhat = P * zhat"); SILC_EXEC("U_L = sparse(U_L-row, U_col, Pzhat, 3*n, s)"); SILC_EXEC("Ump = sparts(U_R-row, U_col, Pzhat, 3*n, s)"); SILC_EXEC("tmp = sqrt(zhat *@ K_stiff *@ L)"); SILC_EXEC("cl = diag(X * sqrt(K_stiff))"); SILC_EXEC("cl = diag(X * sqrt(K_stiff))"); SILC_EXEC("cl = diag(X * tmp); D = diag(tmp)"); SILC_EXEC("Al = Y_LT * cl - Y_RT * cl; Tl = -Al * Al'"); SILC_EXEC("Al = Y_LT * cl - Y_RT * cl; Tl = -Al * Al'"); SILC_EXEC("al = U_L * D - U_R * D; T3 = -A3 * A3'"); SILC_EXEC("DfDx = Tl - T2 + T3");</pre>	
<pre>/* 2. Solve Adv = b */ SILC_EXEC("A = M - (dt * dt) * DfDx - dt * DfDv"); SILC_EXEC("b = dt * (f + dt * (DfDx * v))"); SILC_EXEC("dv = A \\ b");</pre>	
<pre>/* 3. Update particle motion */ SILC_EXEC("v += dv *@ fixed"); SILC_EXEC("x += dt * v");</pre>	













E	xampl	e #3:	Res	ults
$p_{\rm opt}$ is acc	urate or <mark>a</mark>	e <sub>rel</sub> is sm	all (< 10	<mark>)%)</mark> : OK
$\mathcal{E}_{rel}$ is larg	<mark>e</mark> : NG (2	of 16)		
	Number of time steps			
	10	20	30	40
N = 71	2 4 (4) 0%	2 (6) 26.1%	3 (8) 8.6%	3 (8) 8.0%
N = 10	0 <sup>2</sup> 5 (10) 15.4%	5 (8) 7.3%	5 (7) 6.0%	5 (8) 5.7%
N = 14	2 <sup>2</sup> 8 (9) 7.0%	9 (13) 1.9%	8 (11) 1.1%	8 (11) 2.9%
N = 20	0 <sup>2</sup> 12 (12) 0%	12 (15) 0.1%	12 (13) 1.0%	12 (11) 1.0%

Upper:  $p_{opt}$  (the true optimal number of threads in parentheses) Lower: Relative error  $\varepsilon_{rel}$  in the execution time measured with  $p_{op}$ 

#### APT mechanism for SILC

- Outline
  - 1. A server collects *n* samples of execution time with different numbers of threads. For each number of threads, timing is done *m* times and the shortest is picked
  - 2. The server learns a performance model using the least squares method and predicts  $p_{\rm opt}$
  - 3. The server continues the simulation with the optimal number of threads
- Open issue
  - How to determine *n* and *m*

#### Summary

- Proposal of simple performance modeling for time-dependent simulations in SILC
  - Use of the least squares method - Accurate prediction of  $p_{out}$
- · Outline of an APT mechanism for SILC
- Future work
  - Implementation of the APT mechanism
  - How to determine the two parameters n and m

#### Advertisement

· SILC v1.2 is freely available at

#### http://ssi.is.s.u-tokyo.ac.jp/silc/

- Source (Unix/Linux, Windows, Mac OS X)
- Precompiled binary package for Windows
- Documentation, sample programs